

The Excluded 3-minors for V_f -safe Delta-matroids

By

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ABSTRACT. Vf-safe delta-matroids have the desirable property of behaving well under certain duality operations. Several important classes of delta-matroids are known to be vf-safe, including the class of ribbon-graphic delta-matroids, which is related to the class of ribbon graphs or embedded graphs in the same way that graphic matroids correspond to graphs. In this paper, we characterize vf-safe delta-matroids and ribbon-graphic delta-matroids by finding the minimal obstructions, called 3-minors, to belonging to the class. We find the unique (up to twisted duality) excluded 3-minor within the class of set systems for the class of vf-safe delta-matroids. Geelen and Oum [17] found the 166 (up to twists) excluded minors for ribbon-graphic delta-matroids. By translating Bouchet's characterization of circle graphs to the language of 3-minors, we show that this class can also be characterized amongst delta-matroids by a set of three excluded 3-minors up to twisted duality.

1. INTRODUCTION

A *set system* is a pair $S = (E, \mathcal{F})$, where E , or $E(S)$, is a set, called the *ground set*, and \mathcal{F} , or $\mathcal{F}(S)$, is a collection of subsets of E . (All set systems in this paper have finite ground sets.) The members of \mathcal{F} are the *feasible sets*. We say that S is *proper* if $\mathcal{F} \neq \emptyset$.

A matroid M has many associated set systems with $E = E(M)$. The most important of these from the perspective of this paper has $\mathcal{F} = \mathcal{B}(M)$, the set of bases of M . Recall that the bases of a matroid satisfy the following exchange property: for any $B_1, B_2 \in \mathcal{B}(M)$ and for each element $x \in B_1 - B_2$, there is a $y \in B_2 - B_1$ for which $B_1 \Delta \{x, y\} \in \mathcal{B}(M)$. To get the definition of a delta-matroid, replace set differences by symmetric differences. Thus, as introduced by Bouchet in [2], a *delta-matroid* is a proper set system $D = (E, \mathcal{F})$ for which \mathcal{F} satisfies the *delta-matroid symmetric exchange axiom*:

(SE) for all triples (X, Y, u) with X and Y in \mathcal{F} and $u \in X \Delta Y$, there is a $v \in X \Delta Y$ (perhaps u itself) such that $X \Delta \{u, v\}$ is in \mathcal{F} .

Clearly every matroid $(E(M), \mathcal{B}(M))$ is a delta-matroid.

Just as there is a mutually-enriching interplay between matroid theory and graph theory, the theory of delta-matroids has substantial connections with the theory of embedded graphs or equivalently ribbon graphs; see [13, 14]. Brijder and Hooeboom [9, 10, 11] introduced the operation of loop complementation, which we define in the next paragraph. This operation is natural for the class of binary delta-matroids and its subclass of ribbon-graphic delta-matroids. These classes are closed under loop complementation, but that is not true for the class of all delta-matroids. We investigate when loop complementation of a delta-matroid yields a delta-matroid.

For a set system $S = (E, \mathcal{F})$ and $e \in E$, we define $S + e$ to be the set system

$$(1.1) \quad S + e = (E, \mathcal{F} \Delta \{F \cup e : e \notin F \in \mathcal{F}\}).$$

(As in matroid theory, we often omit set braces from singletons.) Note that $(S+e)+e = S$ and that $S+e$ is proper if and only if S is proper. It is straightforward to check that if $e_1, e_2 \in E$ then $(S+e_1)+e_2 = (S+e_2)+e_1$. Thus if $X = \{e_1, \dots, e_n\}$ is a subset of E , then the set system $S+X$ is unambiguously defined by

$$(1.2) \quad S+X = ((S+e_1)+\dots)+e_n.$$

This operation is called the *loop complementation of S on X* . There is a natural operation of embedded graphs, namely *partial Petriality*, to which loop complementation corresponds. More precisely if two embedded graphs are partial Petrials of each other then their ribbon graphic delta-matroids are related by a loop complementation [14, Section 4].

For a delta-matroid D and element $e \in E(D)$, the set system $D+e$ need not be a delta-matroid. Consider, for example, the delta-matroid $D_3 = (\{a, b, c\}, 2^{\{a,b,c\}} - \{\{a, b, c\}\})$. Then $D_3 + \{a, b, c\}$ is the set system $(\{a, b, c\}, \{\emptyset, \{a, b, c\}\})$. This is not a delta-matroid. In fact, it is an excluded minor for the class of delta-matroids [1].

Another operation on delta-matroids is the twist. For $A \subseteq E$, the *twist of S on A* , which is also called the *partial dual of S with respect to A* , denoted $S * A$, is given by

$$S * A = (E, \{F \triangle A : F \in \mathcal{F}\}).$$

Note that $(S * A) * A = S$. The *dual S^** of S is $S * E$. In contrast with loop complementation, each twist of a delta-matroid is a delta-matroid. Apart from the dual, the twists of a matroid $(E(M), \mathcal{B}(M))$ are generally not matroids, as discussed in [14, Theorem 3.4].

Two set systems are said to be *twisted duals* of one another if one may be obtained from the other by a sequence of twists and loop complementations. Following [11], a delta-matroid is said to be *vf-safe* if all of its twisted duals are delta-matroids. (The term vf-safe is short for ‘vertex-flip safe’. Both of the terms vf-safe and loop complementation are named for operations on graphs representing binary delta-matroids [9], which we discuss in Section 5.)

Delta-matroids belonging to certain important classes are known to be vf-safe. In fact, every twisted dual of a ribbon-graphic delta-matroid is a ribbon-graphic delta-matroid [14, Theorem 2.1, Theorem 4.1], and every twisted dual of a binary delta-matroid is a binary delta-matroid [11, Theorem 8.2]. Brijder and Hooeboom showed that quaternary matroids are vf-safe [12], although, as mentioned earlier, matroids are not closed under twists.

In the main result of this paper, Theorem 4.4, we identify D_3 as essentially the unique obstacle for a delta-matroid to be vf-safe. More precisely, we show that the excluded 3-minors for the class of vf-safe delta-matroids within the class of set systems comprise the 28 set systems that are the twisted duals of D_3 . These set systems are given in Tables 2–7. In the final section of the paper, we relate 3-minors to other minor operations that have appeared in the literature, and we apply Theorem 4.4 to recast some known results using short lists of excluded 3-minors.

2. BACKGROUND

Let $S = (E, \mathcal{F})$ be a proper set system. An element $e \in E$ is a *loop* of S if no set in \mathcal{F} contains e . If e is in every set in \mathcal{F} , then e is a *coloop*. If e is not a loop, then the *contraction of e from S* , written S/e , is given by

$$S/e = (E - e, \{F - e : e \in F \in \mathcal{F}\}).$$

If e is not a coloop, then the *deletion of e from S* , written $S \setminus e$, is given by

$$S \setminus e = (E - e, \{F \subseteq E - e : F \in \mathcal{F}\}).$$

If e is a loop or a coloop, then one of S/e and $S \setminus e$ has already been defined, so we can set $S/e = S \setminus e$. Any sequence of deletions and contractions, starting from S , gives another set system S' , called a *minor* of S . Each minor of S is a proper set system.

The order in which elements are deleted or contracted can matter. See [1] for an example. However, for disjoint subsets X and Y of E , if some set in \mathcal{F} is disjoint from X and contains Y , then the deletions and contractions in $S \setminus X/Y$ can be done in any order, and

$$S \setminus X/Y = (E - (X \cup Y), \{F - Y : F \in \mathcal{F} \text{ and } Y \subseteq F \subseteq E - X\}).$$

The following lemma, which is [1, Lemma 2.1], shows that all minors of a proper set system are of this type.

Lemma 2.1. *For any minor S' of a proper set system $S = (E, \mathcal{F})$, there are disjoint subsets X and Y of E such that*

$$S' = S \setminus X/Y = (E - (X \cup Y), \{F - Y : F \in \mathcal{F} \text{ and } Y \subseteq F \subseteq E - X\}).$$

Bouchet and Duchamp [3] showed that, if S is a delta-matroid and $S' = S \setminus X/Y$, then S' is a delta-matroid and S' is independent of the order of the deletions and contractions.

The following definition from [9] is equivalent to that given in equations (1.1)–(1.2). Equivalence can be shown by a routine induction on $|A|$.

Definition 2.2. *If $S = (E, \mathcal{F})$ is a set system and A is a subset of E , then the loop complementation of S by A , denoted by $S + A$, is the set system on E such that F is feasible in $S + A$ if and only if S has an odd number of feasible sets F' with $F - A \subseteq F' \subseteq F$.*

Note that if $A = \{e\}$, then the feasible sets of $S + e$ that do not contain e are the same as those of S , and a set F containing e is feasible in $S + e$ if and only if one but not both of F and $F - e$ is feasible in S . That is, so long as e is not a loop or coloop,

$$\mathcal{F}(S + e) = \mathcal{F}(S \setminus e) \cup \{F \cup e : F \in \mathcal{F}(S \setminus e) \triangle \mathcal{F}(S/e)\}.$$

If e is a loop, then $\mathcal{F}(S + e) = \mathcal{F} \cup \{F \cup e : F \in \mathcal{F}\}$. If e is a coloop, then $S + e = S$.

The twist and loop complementation operations interact in the following way. If A and B are disjoint subsets of E then $(S + A) * B = (S * B) + A$ (a two-element case check and routine induction suffice to verify this), but in general $(S * A) + A \neq (S + A) * A$. However $((S + A) * A) + A = ((S * A) + A) * A$ (see [9]). It follows that there are at most six twisted duals of S with respect to a fixed set A . These relations ensure that any twisted dual of S is of the form $((S * X) + Y) * Z$ for suitably chosen subsets X , Y and Z of E with $X \subseteq Z$. The first relation is used in the proof of the following result.

Lemma 2.3. *Let $S = (E, \mathcal{F})$ be a proper set system, and let a and b be distinct elements of E . Then*

- (i) $S + a \setminus a = S \setminus a$,
- (ii) $S + a \setminus b = S \setminus b + a$, and
- (iii) $S + a/b = S/b + a$.

Proof. If a is a coloop of S , then $S + a = S$, so statement (i) follows. Also, a is not a coloop of S if and only if it is not a coloop of $S + a$, in which case the feasible sets avoiding a are the same in S and $S + a$ by the definition.

For statement (ii), observe that b is a coloop of $S + a$ if and only if it is a coloop of S . When b is not a coloop of S , statement (ii) holds since for each side, the feasible sets are the sets F with $b \notin F$ for which an odd number of the sets X with $F - a \subseteq X \subseteq F$ are in \mathcal{F} . When b is a coloop of S , we need to show that $S + a/b = S/b + a$. This holds since

for each side, the feasible sets are the sets F with $b \notin F$ for which an odd number of the sets X with $(F - a) \cup b \subseteq X \subseteq F \cup b$ are in \mathcal{F} .

It is easy to check that $S'/e = S' * e \setminus e$, so, using statement (ii), we get statement (iii):

$$S + a/b = ((S + a) * b) \setminus b = ((S * b) + a) \setminus b = ((S * b) \setminus b) + a = S/b + a. \quad \square$$

The counterpart, for contractions, of statement (i) is false, as taking $S = D_3$ shows.

3. 3-MINORS

We introduce a third minor operation on set systems. For a proper set system S , we define $S \ddagger e$ to be the set system $(S + e)/e$. This minor operation was first defined by Ellis–Monaghan and Moffatt [15] for ribbon graphs and in an equivalent way by Brijder and Hooeboom [10] for delta-matroids. The notation \ddagger is new, but it seems appropriate to shorten the unwieldy $+e/e$ notation. Motivation for this definition comes from two directions. First, one way to define the Penrose polynomial of a ribbon graph is by specifying a recursive relation analogous to the deletion-contraction recurrence of the chromatic polynomial with this minor operation replacing contraction. For this reason, following a suggestion of Iain Moffatt [18], we propose calling the operation *Penrose contraction*. Second, there is a class of combinatorial objects called multimatroids [6, 7, 8], of which tight 3-matroids are a particular subclass. Brijder and Hooeboom [10] showed that tight 3-matroids are equivalent (in a sense that we do not make precise here) to vf-safe delta-matroids. Tight 3-matroids have three minor operations corresponding to deletion, contraction, and Penrose contraction in vf-safe delta-matroids.

Any sequence of the three minor operations, starting from S , gives another proper set system S' , called a *3-minor* of S . A collection \mathcal{C} of proper set systems is *3-minor closed* if every 3-minor of every member of \mathcal{C} is in \mathcal{C} . Given such a collection \mathcal{C} , a proper set system S is an *excluded 3-minor* for \mathcal{C} if $S \notin \mathcal{C}$ and all other 3-minors of S are in \mathcal{C} . A proper set system belongs to \mathcal{C} if and only if none of its 3-minors is an excluded 3-minor for \mathcal{C} . Thus, the excluded 3-minors determine \mathcal{C} ; they are the 3-minor-minimal obstructions to membership in \mathcal{C} .

For a given class \mathcal{C} , there may be substantially fewer excluded 3-minors than excluded minors. For instance, in [17], Geelen and Oum found 166 delta-matroids that, up to twists, are the excluded minors for ribbon-graphic delta-matroids within the class of binary delta-matroids. In contrast, in Theorem 5.8, we show that every binary matroid that does not have a twisted dual of one of three delta-matroids as a 3-minor is ribbon-graphic.

An element e is called a *pseudo-loop* of S if e is a loop of $S + e$. The next lemma characterizes these elements.

Lemma 3.1. *For an element e in a proper set system S , the following statements are equivalent:*

- (i) e is a loop of $S + e$, that is, a pseudo-loop of S ,
- (ii) $F \cup e \in \mathcal{F}(S)$ if and only if $F \in \mathcal{F}(S)$, and
- (iii) $S * e = S$.

Pseudo-loops of S are neither loops nor coloops of S . Furthermore, $S \ddagger e = S \setminus e = S/e$ if and only if e is a loop, a coloop, or a pseudo-loop of S .

Proof. The equivalence of statements (i)–(iii) is immediate from the definitions. Statement (ii) implies that pseudo-loops are neither loops nor coloops. If e is a loop of S , then $S \ddagger e = S \setminus e$ since $\mathcal{F}(S + e) = \mathcal{F}(S) \cup \{F \cup e : F \in \mathcal{F}(S)\}$; also, $S \setminus e = S/e$ by definition. If e is a coloop of S , then $S \ddagger e = S/e$ since $S + e = S$; also, $S \setminus e = S/e$ by

definition. If e is a pseudo-loop of S , then statements (i) and (ii) gives the equality. If e is not a loop, a coloop, or a pseudo-loop of S , then $S \setminus e \neq S/e$ by the failure of statement (ii) and the fact that some, but not all, sets in $\mathcal{F}(S)$ contain e . \square

The following two results show that one may choose the operations used to form a 3-minor in such a way that they commute.

Lemma 3.2. *Let $S = (E, \mathcal{F})$ be a proper set system, and let X, Y , and Z be pairwise disjoint subsets of E . If there is a set F with*

$$(3.1) \quad F \subseteq E - (X \cup Y \cup Z) \quad \text{and} \quad |\mathcal{F} \cap \{F' : F \cup Y \subseteq F' \subseteq F \cup Y \cup Z\}| \text{ is odd,}$$

then the minor operations in $S \setminus X/Y \ddagger Z$ can be done in any order and a set F is feasible in $S \setminus X/Y \ddagger Z$ if and only if it satisfies Condition (3.1).

Proof. A set F meets Condition (3.1) if and only if $F \subseteq E - (X \cup Y \cup Z)$ and $F \cup Y \cup Z$ is in $\mathcal{F}(S + Z)$. If there is at least one set satisfying Condition (3.1), the remarks preceding Lemma 2.1 imply that the deletions and contractions in forming $(S + Z) \setminus X/(Y \cup Z)$ from $S + Z$ may be done in any order and a set F is feasible in $(S + Z) \setminus X/(Y \cup Z)$ if and only if it satisfies Condition (3.1). Lemma 2.3 implies that we may defer taking a loop complementation of an element in Z until just before it is contracted. The result follows. \square

We next show that for every 3-minor of a proper set system, there are pairwise disjoint sets X, Y and Z satisfying Condition (3.1).

Lemma 3.3. *Let S' be a 3-minor of a proper set system $S = (E, \mathcal{F})$. Then there are pairwise disjoint subsets X, Y , and Z of E such that $S' = S \setminus X/Y \ddagger Z$ and there is a set F satisfying Condition (3.1).*

Proof. Suppose we get S' from S by, for each of e_1, e_2, \dots, e_k in turn, performing one of the three minor operations, giving the sequence of minors $S_0 = S, S_1, \dots, S_k = S'$. Let X be the set of elements e_i in $\{e_1, \dots, e_k\}$ that satisfy at least one of the following conditions:

- (1) e_i is a loop or a pseudo-loop of S_{i-1} , so $S_i = S_{i-1} \setminus e_i$, or
- (2) e_i is not a coloop of S_{i-1} and $S_i = S_{i-1} \setminus e_i$.

Let Y be the set of elements e_i in $\{e_1, \dots, e_k\} - X$ such that e_i is either a coloop of S_{i-1} or $S_i = S_{i-1}/e_i$. Note that if $e_i \in Y$ then it is not a loop in S_{i-1} . Finally let $Z = \{e_1, \dots, e_k\} - (X \cup Y)$, so that Z comprises the elements e_i in $\{e_1, \dots, e_k\}$ for which $S_i = S_{i-1} \ddagger e_i$ but e_i is not a loop, pseudo-loop or coloop. Then there is always at least one set F satisfying Condition (3.1). \square

Table 1 shows the result of applying one of the three minor operations that remove e after taking one of the six twisted duals, with respect to e , of a proper set system. If instead the minor operation removes a different element from that used for the twisted dual, then these operations commute.

We next show that any 3-minor of a twisted dual of a proper set system S is a twisted dual of some 3-minor of S . It is easy to see that the converse is also true.

Lemma 3.4. *Suppose S is a proper set system and S' is a twisted dual of S . If S'' is a 3-minor of S' , then S has a 3-minor that is a twisted dual of S'' .*

Proof. There are subsets A and B of $E(S)$ such that we obtain S'' from S by first forming a twisted dual for each element of A and then performing one of the three minor operations for each element of B . The remarks before this lemma imply that one may reorder these

	/e	\e	‡e
S	S/e	$S \setminus e$	$S \ddagger e$
$S * e$	$S \setminus e$	S/e	$S \ddagger e$
$S + e$	$S \ddagger e$	$S \setminus e$	S/e
$(S + e) * e$	$S \setminus e$	$S \ddagger e$	S/e
$(S * e) + e$	$S \ddagger e$	S/e	$S \setminus e$
$((S * e) + e) * e$	S/e	$S \ddagger e$	$S \setminus e$

TABLE 1. Interaction of minor operations and twisted duality.

operations to first deal with the elements of $A \cap B$, one by one, forming a twisted dual for an element and then a 3-minor before moving on to the next element. According to Table 1 each of these pairs of operations may be replaced by a single 3-minor operation. Next a 3-minor is formed for each element of $B - A$. The resulting set system is a twisted dual of S'' with respect to the elements of $A - B$. \square

4. CHARACTERIZATIONS BY EXCLUDED 3-MINORS

Brijder and Hoogeboom [11] showed that the class of vf-safe delta-matroids is minor-closed. A computer search for excluded minors for this class turns up many examples with apparently little structure. The class of vf-safe delta-matroids is defined using both the twist and loop complementation operations, so it is natural to try to characterize this class using 3-minors. By Lemma 4.1 below, the class of vf-safe delta-matroids is closed under Penrose contraction, so, with the result in [11], it is closed under 3-minors. The main result of this section, Theorem 4.4, gives the excluded 3-minors for the class of vf-safe delta-matroids within the class of set systems.

Lemma 4.1. *If S is vf-safe and $e \in E(S)$, then $S \ddagger e$ is vf-safe.*

Proof. If S is vf-safe, then all of its twisted duals are vf-safe by definition, so $S + e$ is vf-safe. Theorem 8.3 in [11] states that every minor of a vf-safe delta-matroid is vf-safe. Thus $S \ddagger e = S + e/e$ is vf-safe. \square

Let

$$S_i = (\{e_1, e_2, \dots, e_i\}, \{\emptyset, \{e_1, e_2, \dots, e_i\}\}).$$

Let \mathcal{S} be the set of all twists of the set systems in $\{S_3, S_4, \dots\}$. Let

- $T_1 = (\{a, b, c\}, \{\emptyset, \{a, b\}, \{a, b, c\}\});$
- $T_2 = (\{a, b, c\}, \{\emptyset, \{a, b\}, \{a, c\}, \{a, b, c\}\});$
- $T_3 = (\{a, b, c\}, \{\emptyset, \{a\}, \{a, b\}, \{a, b, c\}\});$
- $T_4 = (\{a, b, c\}, \{\emptyset, \{a\}, \{a, b\}, \{a, c\}, \{a, b, c\}\});$
- $T_5 = (\{a, b, c, d\}, \{\emptyset, \{a, b\}, \{a, b, c, d\}\});$
- $T_6 = (\{a, b, c, d\}, \{\emptyset, \{a, b\}, \{a, c\}, \{a, b, c, d\}\});$
- $T_7 = (\{a, b, c, d\}, \{\emptyset, \{a, b\}, \{a, c\}, \{a, d\}, \{a, b, c, d\}\});$
- $T_8 = (\{a, b, c, d\}, \{\emptyset, \{a\}, \{a, b\}, \{a, c\}, \{a, d\}, \{a, b, c, d\}\}).$

Let \mathcal{T} be the set of all twists of the set systems in $\{T_1, T_2, \dots, T_8\}$. By the following result from [1, Theorem 5.1], these are all of the excluded minors for delta-matroids within the class of set systems.

Theorem 4.2. *A proper set system S is a delta-matroid if and only if S has no minor isomorphic to a set system in $\mathcal{S} \cup \mathcal{T}$.*

The following lemma is key for finding the excluded 3-minors for vf-safe delta-matroids within the class of set systems.

Lemma 4.3. *Let S be an excluded 3-minor for the class of vf-safe delta-matroids. Then S has a twisted dual that is isomorphic to a set system in $\mathcal{S} \cup \mathcal{T}$.*

Proof. Such an excluded 3-minor S either is not a delta-matroid and all of its other minors are delta-matroids, or it is a delta-matroid and has a twisted dual S' that is not a delta-matroid. In the former case S is isomorphic to a set system in $\mathcal{S} \cup \mathcal{T}$ and the lemma holds. In the latter case S' has a minor S'' isomorphic to a member of $\mathcal{S} \cup \mathcal{T}$. By Lemma 3.4, S has a 3-minor S''' that is a twisted dual of S'' . Therefore S''' is not a vf-safe delta-matroid. The only 3-minor of S that is not a vf-safe delta-matroid is S itself. Hence $S = S'''$ and the lemma holds. \square

To connect the next result with the remarks in Section 1, note that $D_3 + \{a, b, c\} = S_3$.

Theorem 4.4. *A proper set system is a vf-safe delta-matroid if and only if it has no 3-minor that is isomorphic to a twisted dual of S_3 .*

Proof. All proper set systems with two elements are delta-matroids, and therefore each one is vf-safe, so the twisted duals of S_3 are excluded 3-minors for the class of vf-safe delta-matroids. By Lemma 4.3 every excluded 3-minor for the class of vf-safe delta-matroids must be a twisted dual of a set system in $\mathcal{S} \cup \mathcal{T}$. We first consider the set systems with three-element ground sets. We have $T_1^* + c = S_3$ and $T_2^* + \{b, c\} \simeq T_3 + a = T_1$ and $T_4 + a = T_2$, so every excluded 3-minor of size three is a twisted dual of S_3 .

Lastly, we show that no other set system in $\mathcal{S} \cup \mathcal{T}$ is an excluded 3-minor. Lemma 3.4 implies that each twisted dual of an excluded 3-minor is an excluded 3-minor, so it suffices to show that each of T_5, T_6, T_7, T_8 , and S_n , for $n \geq 4$, has a smaller member of $\mathcal{S} \cup \mathcal{T}$ as a 3-minor. Indeed, $S_n \ddagger e_n = S_{n-1}$, for $n \geq 4$, $T_5 \ddagger d = T_1$, $T_6 \ddagger d = T_8 \ddagger d = T_2$, and $T_7 \ddagger d = T_4$. \square

As stated in the introduction, there are 28 twisted duals of S_3 , up to isomorphism. These excluded 3-minors are listed in Tables 2–7.

5. 3-MINORS AND VERTEX MINORS

We now explain the link between 3-minors and vertex minors in binary delta-matroids. Vertex minors are well-studied, but are only defined for binary delta-matroids. In contrast, the concept of a 3-minor is relatively unstudied, but is important due to its direct correlation with ribbon-graph operations and its applicability beyond binary delta-matroids. For this reason, and for completeness, we give a full explanation here. Although the key ideas presented here are not new, the link between vertex minors and 3-minors has not previously been fully described.

A delta-matroid is *normal* if the empty set is feasible. A delta-matroid is *even* if for every pair F_1 and F_2 of its feasible sets $|F_1 \triangle F_2|$ is even. Equivalently, the sizes of all its feasible sets are of the same parity. Let M denote a symmetric binary matrix with rows and columns indexed by $[n] = \{1, \dots, n\}$. Take $E = [n]$ and \mathcal{F} to be the collection of subsets S of $[n]$ for which the principal submatrix of M comprising the rows and columns indexed by elements of S is non-singular. Bouchet [3] showed that $D(M) = (E, \mathcal{F})$ is a delta-matroid. We call such delta-matroids *basic binary*. (In the literature, what we have called basic binary delta-matroids are often called graphic, but we prefer to avoid this term to prevent confusion with ribbon-graphic delta-matroids.) A delta-matroid is *binary* [3] if it is a twist of a basic binary delta-matroid.

It follows immediately from the definition that every basic binary delta-matroid is normal and that a basic binary delta-matroid is uniquely determined by its feasible sets of size at most two. A well-known result of linear algebra states that a symmetric matrix with an odd number of rows (and columns) and zero diagonal is singular. Consequently a basic binary delta-matroid is even if and only if it has no feasible sets of size one.

Let A be a matrix over an arbitrary field with rows and columns indexed by $[n]$, and let X be a subset of $[n]$ such that the principal sub-matrix $P = A[X]$ is non-singular. Suppose without loss of generality that $A = \begin{pmatrix} P & Q \\ R & S \end{pmatrix}$. Then the matrix $A * X$ is defined by

$$A * X = \begin{pmatrix} P^{-1} & -P^{-1}Q \\ RP^{-1} & S - RP^{-1}Q \end{pmatrix}.$$

Note that if A is a symmetric binary matrix then $A * X$ is symmetric. The following result is due to Tucker [20].

Theorem 5.1. *Let A be a matrix over an arbitrary field with rows and columns indexed by $[n]$, and let X be a subset of $[n]$ such that the principal sub-matrix $P = A[X]$ is non-singular. Then for every subset Y of $[n]$, we have*

$$\det((A * X)[Y]) = \frac{\det(A[X \Delta Y])}{\det(A[X])}.$$

In particular for any subset Y of $[n]$, the principal submatrix $(A * X)[Y]$ is non-singular if and only if the principal submatrix $A[X \Delta Y]$ is non-singular.

The following corollary is immediate.

Corollary 5.2. *Suppose that A is a binary matrix, and X is a feasible set of $D(A)$. Then $D(A) * X = D(A * X)$.*

See [3] for an alternative proof of this result that holds for arbitrary fields. A consequence of this corollary is that every normal twist of a basic binary delta-matroid is basic binary.

A *looped simple graph* is a graph without parallel edges but in which each vertex is allowed to have up to one loop. Much of the time we will forbid loops; we call such graphs *loopless simple graphs*. Recall that basic binary delta-matroids are completely determined by their feasible sets with size two or fewer. Clearly basic binary delta-matroids on the set $[n]$ are in one-to-one correspondence with looped simple graphs with vertex set $[n]$; likewise, even basic binary delta-matroids on $[n]$ are in one-to-one correspondence with loopless simple graphs with vertex set $[n]$.

We take adjacency matrices to always be binary. Given a looped simple graph G and its adjacency matrix A , we let $D(G)$ denote the basic binary delta-matroid $D(A)$. If X is a subset of the edges of G , then X labels a subset of the rows and columns of A , and we define $G * X$ to be the looped simple graph with adjacency matrix $A * X$.

We now consider various transformations that may be applied to G and their effect on $D(G)$.

The loop complementation operation of Brijder and Hoogeboom was first defined in terms of basic binary delta-matroids. If G is a looped simple graph and v is a vertex of G , then the loop complementation $G + v$ is formed by toggling the loop at v , that is, removing a loop if there is one at v and adding one at v if there is no loop there.

The following lemma from [9] is straightforward.

Lemma 5.3. *Let G be a looped simple graph with vertex v . Then $D(G + v) = D(G) + v$.*

Our next operation is local complementation. There are several variations in the definition of local complementation: see, for instance, [19]. We will only require certain cases of what is defined there. For a looped simple graph G with vertex v , let $N_G(v)$ denote the *open neighbourhood* of v , that is, the set of vertices, excluding v , that are adjacent to v in G . We will generally write N instead of N_G when there is no possibility of confusion. The *local complementation* of G at v , denoted by G^v , is formed by toggling the adjacencies between pairs of neighbours of v , that is, for every distinct pair x, y from the neighbourhood of v , delete edge xy if x and y are adjacent in G and add edge xy if x and y are not adjacent in G . Additionally, if there is a loop at v , then the loop status of every vertex in the open neighbourhood of v is toggled. In both cases, adjacencies involving one or more non-neighbours of v or v itself are unchanged and the presence or not of a loop at v is unaffected. To distinguish the two cases, it will be helpful to refer to local complementation at v as *simple local complementation* when v is loopless, and *non-simple local complementation* when there is a loop at v .

For delta-matroid D and subset $A \subseteq E(D)$, let $D\bar{*}A$ denote the *dual pivot on A* , which is equal to $D + A * A + A$. The following result is implicit in the results of [19], but is not expressed in this form.

Proposition 5.4. *Let G be a loopless simple graph with vertex v . Then $D(G^v) = (D(G)\bar{*}v) + N(v)$.*

Proof. Let A be the adjacency matrix of G . Then A is symmetric and all of its diagonal entries are zero. Notice that the simple local complementation G^v can be formed by adding a loop at v , performing the non-simple local complementation at v and then removing the loops added at v and all of its neighbours.

We have $D(G + v) = D(G) + v$. Assume without loss of generality that $v = 1$ and let $Z = [n] - 1$. Then the adjacency matrix of $G + v$ is $\begin{pmatrix} 1 & c \\ c^t & A[Z] \end{pmatrix}$ for some vector c . Then it follows from Corollary 5.2 that $(D(G) + v) * v = D((G + v) * v) = D(A')$ where $A' = \begin{pmatrix} 1 & c \\ c^t & A[Z] + c^t c \end{pmatrix}$.

A diagonal entry of $c^t c$ is non-zero if it corresponds to a neighbour of v and an off-diagonal entry of $c^t c$ is non-zero if both the row and column correspond to neighbours of v . Thus $(D(G) + v) * v = D(G')$ where G' is formed from G by adding a loop at v and performing the non-simple local complementation at v . Now G' has loops at v and at all neighbours of v , so

$$D(G^v) = D(G' + v + N(v)) = D(G') + v + N(v) = (D(G)\bar{*}v) + N(v). \quad \square$$

The corollary below is well-known and follows from the previous result.

Corollary 5.5. *Let G be a loopless simple graph with adjacent vertices v and w . Then $D(((G^v)^w)^v) = D(G) * \{v, w\}$.*

Proof. We have

$$D(((G^v)^w)^v) = ((D(G)\bar{*}v + N(v))\bar{*}w + N_{G^v}(w))\bar{*}v + N_{(G^v)^w}(v).$$

It follows from the discussion before Lemma 2.3 that one may reorder the loop complement and twist operations so that those involving a particular vertex of G are done consecutively. The result follows by considering the effect of the operations involving each vertex of G separately and noting that

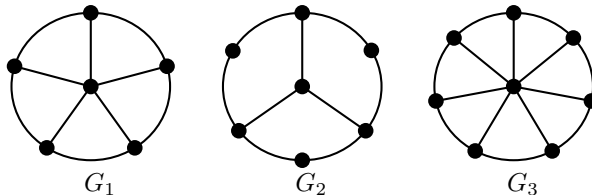


FIGURE 1. A complete set of circle graph obstructions.

- (1) a common neighbour of v and w in G is a neighbour of v but not w in both G^v and $(G^v)^w$,
- (2) a vertex other than w that is a neighbour of v but not w in G is a neighbour of both v and w in G^v and of w but not v in $(G^v)^w$, and
- (3) a vertex other than v that is a neighbour of w but not v in G is a neighbour of both v and w in $(G^v)^w$ and of w but not v in G^v . \square

A *vertex minor* of a looped simple graph G is formed from G by a sequence of local complementations and deletions of vertices. It is easy to check that if v and w are different vertices of an unlooped simple graph, then $(G^v) \setminus w = (G \setminus w)^v$ and thus one may assume that all the local complementations are done first.

Perhaps the most important use of vertex minors is Bouchet's characterization of circle graphs. A *chord diagram* is a collection of chords of a circle. Chord diagrams are in one-to-one correspondence with orientable ribbon graphs with one vertex. (For more information on ribbon graphs, see [16] or [14].) To see this think of the circle and its interior as the vertex of a ribbon graph and for each chord attach a ribbon to the vertex at the points corresponding to the endpoints of the chord. Clearly two chords intersect if and only if the corresponding ribbons e_1 and e_2 are interlaced, that is, as one traverses the vertex one meets an end of e_1 , then an end of e_2 , then the other end of e_1 , and finally the other end of e_2 . An unlooped simple graph is a *circle graph* if it is the intersection graph of the chords in a chord diagram, that is, there is a vertex corresponding to each chord and they are adjacent if and only if the chords cross. Equivalently a circle graph is the interlacement graph of an orientable ribbon graph with one vertex: it has a vertex for each ribbon and two vertices are adjacent if the corresponding ribbons are interlaced. Bouchet established the following result [5].

Theorem 5.6. *An unlooped simple graph is a circle graph if and only if it has no vertex minor isomorphic to the graphs G_1 , G_2 or G_3 depicted in Figure 1.*

We are now ready to state the link between 3-minors and vertex minors.

- Theorem 5.7.** (1) *Let G be a unlooped simple graph and let H be a vertex minor of G . Then $D(H)$ is a 3-minor of $D(G)$.*
- (2) *Let D be a twisted dual of a basic binary delta-matroid and let D' be a 3-minor of D . Then there are graphs G and G' such that $D(G)$ and $D(G')$ are twisted duals of D and D' respectively, and G' is a vertex minor of G .*

Proof. For part (1), note that a vertex minor of an unlooped simple graph is obtained by a sequence of local complementations and vertex deletions. The result follows from Proposition 5.4 and the fact that if v is a vertex of G then $D(G \setminus v) = D(G) \setminus v$.

For part (2), let F be a feasible set of D and let

$$B = \{e \in E(D) : \{e\} \in \mathcal{F}(D * F)\}.$$

The remark following Corollary 5.2 implies that $D * F$ is basic binary, so $(D * F) + B$ is an even basic binary delta-matroid, so $(D * F) + B = D(G)$ for some unlooped simple graph G . It follows from Lemma 3.4 that there is a delta-matroid D'' that is a 3-minor of $D(G)$ and a twisted dual of D' . We shall prove by induction on k that if G is an unlooped simple graph and D'' is a 3-minor of $D(G)$ with k fewer elements, then there is an unlooped simple graph G' that is a vertex minor of G and such that $D(G')$ is a twisted dual of D'' . The result then follows.

If $k = 0$, then take $G' = G$. Otherwise D'' is obtained from $D(G)$ by a sequence of k deletions, contractions and Penrose contractions. Suppose that the first operation is the deletion of v . Then take $G'' = G \setminus v$, which is a vertex minor of G . Furthermore $D(G) \setminus v = D(G'')$ and D'' is a 3-minor of $D(G'')$ with $k - 1$ fewer edges. Hence, by induction, there is an unlooped simple graph G' that is a vertex minor of G'' and hence of G , and such that $D(G')$ is a twisted dual of D'' . Suppose next that the first operation is the Penrose contraction of v . Then take $G'' = (G^v) \setminus v$. We have

$$\begin{aligned} D(G'') &= D(G^v \setminus v) \\ &= (((D(G) + v) * v) + v) + N(v) \setminus v \\ &= (((D(G) * v) + v) * v) \setminus v + N(v) \\ &= (((D(G) * v) + v)/v) + N(v) \\ &= (D(G) \ddagger v) + N(v). \end{aligned}$$

(The last equality uses Table 1.) Now $D(G'')$ is a twisted dual of $D(G) \ddagger v$, so it has a 3-minor D''' with $k - 1$ fewer elements that is a twisted dual of D'' . Hence, by induction, there is an unlooped simple graph G' that is a vertex minor of G'' such that $D(G')$ is a twisted dual of D''' and consequently of D'' . In the final case the first operation is the contraction of v . If v is an isolated vertex of G then v appears in no feasible set of $D(G)$ of size at most two and consequently in no feasible set of $D(G)$ of any size. Thus v is a loop of $D(G)$ and $D(G)/v = D(G) \setminus v = D(G \setminus v)$. If v is not an isolated vertex of G then let w be a neighbour of v . We have

$$\begin{aligned} D(((G^v)^w)^v \setminus v) &= D(((G^v)^w)^v) \setminus v \\ &= (D(G) * \{v, w\}) \setminus v \\ &= (D(G)/v) * w. \end{aligned}$$

The proof of this case is completed in a similar way to the case of Penrose contraction. \square

From the preceding result we obtain the following restatement of Bouchet's result, determining the three binary delta-matroids that are the excluded 3-minors for ribbon-graphic delta-matroids.

Theorem 5.8. *A binary delta-matroid is ribbon-graphic if and only if it has no 3-minor that is a twisted dual of $D(G_1)$, $D(G_2)$ or $D(G_3)$.*

Proof. If D is a binary delta-matroid and v is an element of D then D is ribbon-graphic if and only if $D + v$ is ribbon graphic, because it follows from Theorem 4.1 of [14] that if R is a ribbon graph with $D = D(R)$ then $D + v$ is the delta-matroid corresponding to the ribbon graph formed from R by applying a half-twist to v . Let

$$B = \{e \in E(D) : \{e\} \in \mathcal{F}(D)\}.$$

Then $D + B$ is even and, furthermore, $D + B$ is ribbon-graphic if and only if D is ribbon-graphic. Now $D + B = D(G)$ where G is an unlooped simple graph. Bouchet's Theorem 5.6 states that G is a circle graph if and only if G has no vertex minor isomorphic to G_1, G_2 or G_3 . Equivalently $D + B$ is ribbon-graphic if and only if it has no 3-minor that is a twisted dual of $D(G_1), D(G_2)$ or $D(G_3)$. As $D + B$ is a twisted dual of D , the result follows. \square

We close by presenting excluded 3-minor results for the classes of binary delta-matroids and ribbon graphic delta-matroids that follow from Theorem 4.4. Bouchet [4] proved that every minor of a binary delta-matroid is binary and gave the following excluded-minor characterization of binary delta-matroids.

Theorem 5.9. *A delta-matroid is binary if and only if it does not have a minor isomorphic to any of the following five delta-matroids or their twists.*

- (1) $B_1 = (\{a, b, c\}, \{\emptyset, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\})$;
- (2) $B_2 = S_3 + \{a, b, c\}$;
- (3) $B_3 = (\{a, b, c\}, \{\emptyset, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{a, b, c\}\})$;
- (4) $B_4 = (\{a, b, c, d\}, \{\emptyset, \{a, b\}, \{a, c\}, \{a, d\}, \{b, c\}, \{b, d\}, \{c, d\}\})$;
- (5) $B_5 = (\{a, b, c, d\}, \{\emptyset, \{a, b\}, \{a, d\}, \{b, c\}, \{c, d\}, \{a, b, c, d\}\})$.

We obtain corollaries of this result characterizing binary delta-matroids in terms excluded 3-minors.

Corollary 5.10. *A vf-safe delta-matroid is binary if and only if it has no 3-minor that is a twisted dual of B_1 .*

Proof. Theorem 8.2 of [11] states that every twisted dual of a binary delta-matroid is a binary delta-matroid. In particular every binary delta-matroid is vf-safe. Moreover, every 3-minor of a binary delta-matroid is binary. The delta-matroid B_1 is vf-safe and all of its 3-minors are binary. Thus all of its twisted duals are excluded 3-minors for the class of binary delta-matroids.

Suppose that D is a vf-safe delta-matroid that is not binary. Then Theorem 5.9 implies that D has a minor isomorphic to a twist of B_i for $1 \leq i \leq 5$. The delta-matroid B_2 is not vf-safe and $B_4 \ddagger d = B_2$, so D has no minor isomorphic to a twist of B_2 or of B_4 . Furthermore $(B_3 + a)^* = B_1$, and $B_5 \ddagger d \simeq B_3$. Thus D has a 3-minor that is isomorphic to a twisted dual of B_1 . \square

By combining this result with Theorem 4.4, we obtain the following.

Corollary 5.11. *A proper set system is a binary delta-matroid if and only if it has no 3-minor that is a twisted dual of B_1 or S_3 .*

Finally we combine the last two results with Theorem 5.8.

Corollary 5.12. *A proper set system is a ribbon graphic delta-matroid if and only if it has no 3-minor that is a twisted dual of $B_1, S_3, D(G_1), D(G_2)$ or $D(G_3)$.*

6. APPENDIX: THE TWISTED DUALS OF S_3

As proved in Theorem 4.4, these twisted duals of S_3 are the excluded 3-minors for vf-safe delta-matroids.

$$S_3 \begin{array}{|c|c|} \hline \emptyset & \{a, b, c\} \\ \hline \end{array}$$

$$S_3 * \{a\} \begin{array}{|c|c|} \hline \{a\} & \{b, c\} \\ \hline \end{array}$$

TABLE 2. All twists of S_3 up to isomorphism.

\emptyset	$\{a\}$	$\{a, b, c\}$	\emptyset	$\{b, c\}$	$\{a, b, c\}$
$S_3 + \{a\}$			$(S_3 + \{a\})^*$		
\emptyset	$\{a\}$	$\{b, c\}$	$\{a\}$	$\{b, c\}$	$\{a, b, c\}$
$(S_3 + \{a\}) * \{a\}$			$(S_3 + \{a\}) * \{b, c\}$		
$\{b\}$	$\{a, b\}$	$\{a, c\}$	$\{b\}$	$\{a, c\}$	
$(S_3 + \{a\}) * \{b\}$			$(S_3 + \{a\}) * \{a, c\}$		

TABLE 3. All twists of $S_3 + \{a\}$ up to isomorphism. Dual pairs are side by side.

\emptyset	$\{a\}$	$\{a, b\}$	$\{a, b, c\}$	\emptyset	$\{c\}$	$\{a, c\}$	$\{a, b, c\}$
$S_3 + \{a, b\}$				$(S_3 + \{a, b\})^*$			
\emptyset	$\{a\}$	$\{a, b\}$	$\{a, b, c\}$	$\{a\}$	$\{a, c\}$	$\{a, b, c\}$	
$(S_3 + \{a, b\}) * \{a\}$				$(S_3 + \{a, b\}) * \{b, c\}$			
$\{c\}$	$\{a, b\}$	$\{a, c\}$	$\{a, b, c\}$	\emptyset	$\{a\}$	$\{a, b\}$	
$(S_3 + \{a, b\}) * \{c\}$				$(S_3 + \{a, b\}) * \{a, b\}$			

TABLE 4. All twists of $S_3 + \{a, b\}$ up to isomorphism. Dual pairs are side by side.

\emptyset	$\{a\}$	$\{a, b\}$	$\{a, b, c\}$	$\{a\}$	$\{a, b\}$	$\{a, b, c\}$	
$S_3 + \{a, b, c\}$				$(S_3 + \{a, b, c\})^*$			
\emptyset	$\{a\}$	$\{a, b\}$	$\{a, b, c\}$	$\{b\}$	$\{a, c\}$	$\{a, b, c\}$	
$S_3 + \{a, b, c\} * \{a\}$				$S_3 + \{a, b, c\} * \{b, c\}$			
$\{c\}$	$\{a, b\}$	$\{a, c\}$	$\{a, b, c\}$	$\{c\}$	$\{b, c\}$		

TABLE 5. All twists of $S_3 + \{a, b, c\}$ up to isomorphism. Dual pairs are side by side.

$\begin{array}{l} \{a\} \\ \{a, b\} \\ \{b, c\} \end{array}$	$\begin{array}{l} \{a, b\} \\ \{a, b, c\} \end{array}$	\emptyset	$\begin{array}{l} \{a\} \\ \{c\} \end{array}$	$\{b, c\}$				
$(S_3 * \{a\}) + \{a, b\}$		$((S_3 * \{a\}) + \{a, b\})^*$						
<table style="border: 1px solid black; margin: auto; padding: 5px;"> <tr> <td style="padding: 2px;">\emptyset</td> <td style="padding: 2px;">$\{b\}$</td> <td style="padding: 2px;">$\{b, c\}$</td> <td style="padding: 2px;">$\{a, b, c\}$</td> </tr> </table>					\emptyset	$\{b\}$	$\{b, c\}$	$\{a, b, c\}$
\emptyset	$\{b\}$	$\{b, c\}$	$\{a, b, c\}$					
$((S_3 * \{a\}) + \{a, b\}) * \{a\}$								
<table style="border: 1px solid black; margin: auto; padding: 5px;"> <tr> <td style="padding: 2px;">$\{a\}$</td> <td style="padding: 2px;">$\{a, b\}$</td> </tr> <tr> <td style="padding: 2px;">$\{c\}$</td> <td style="padding: 2px;">$\{a, c\}$</td> </tr> </table>					$\{a\}$	$\{a, b\}$	$\{c\}$	$\{a, c\}$
$\{a\}$	$\{a, b\}$							
$\{c\}$	$\{a, c\}$							
$((S_3 * \{a\}) + \{a, b\}) * \{b\}$								

TABLE 6. All twists of $(S_3 * \{a\}) + \{a, b\}$ up to isomorphism. Dual pairs are side by side.

$\begin{array}{l} \{a\} \\ \{a, c\} \\ \{b, c\} \end{array}$	$\begin{array}{l} \{a, b\} \\ \{a, c\} \\ \{b, c\} \end{array}$	$\begin{array}{l} \{a\} \\ \{b\} \\ \{c\} \end{array}$	$\{b, c\}$						
$(S_3 * \{a\}) + \{a, b, c\}$		$((S_3 * \{a\}) + \{a, b, c\})^*$							
<table style="border: 1px solid black; margin: auto; padding: 5px;"> <tr> <td style="padding: 2px;">\emptyset</td> <td style="padding: 2px;">$\{b\}$</td> <td style="padding: 2px;">$\{a, b, c\}$</td> </tr> <tr> <td style="padding: 2px;">$\{c\}$</td> <td></td> <td></td> </tr> </table>				\emptyset	$\{b\}$	$\{a, b, c\}$	$\{c\}$		
\emptyset	$\{b\}$	$\{a, b, c\}$							
$\{c\}$									
$((S_3 * \{a\}) + \{a, b, c\}) * \{a\}$		$((S_3 * \{a\}) + \{a, b, c\}) * \{b, c\}$							
<table style="border: 1px solid black; margin: auto; padding: 5px;"> <tr> <td style="padding: 2px;">$\{a\}$</td> <td style="padding: 2px;">$\{a, b\}$</td> <td style="padding: 2px;">$\{a, b, c\}$</td> </tr> <tr> <td style="padding: 2px;">$\{c\}$</td> <td></td> <td></td> </tr> </table>				$\{a\}$	$\{a, b\}$	$\{a, b, c\}$	$\{c\}$		
$\{a\}$	$\{a, b\}$	$\{a, b, c\}$							
$\{c\}$									
$((S_3 * \{a\}) + \{a, b, c\}) * \{b\}$		$((S_3 * \{a\}) + \{a, b, c\}) * \{a, c\}$							

TABLE 7. All twists of $(S_3 * \{a\}) + \{a, b, c\}$ up to isomorphism. Dual pairs are side by side.

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